

The Quest Towards Analytical Solutions of Linked Fault Tree Models using Binary Decision Diagrams^(*)

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- Motivation and issues with current PSA softwares
- Binary Decision Diagrams (BDD) as an alternative
- Research and development at KKL and ETH Zurich
- Insights and outlook
- Presentation of the *NeuralSpectrum Software*



(*) From the PhD thesis:

“Analytical Solutions of Linked Fault Tree Probabilistic Risk Assessments using Binary Decision Diagrams with Emphasis on Nuclear Safety Applications”



Motivation and issues with current PSA softwares

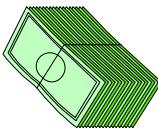
- The (Swiss) NPP have to submit **best-estimate, plant-specific PSA models to the regulatory authority**
 - ☼ Level 1 PSA: Calculation of the Core Damage Frequency (CDF)
 - ☼ Level 2 PSA: Calculation of the Plant Damage State (PDS)
frequencies and associated radiological
consequences
 - ☼ For internal events, area events and external events
 - ☼ Typically, the PSA modeling techniques are based on the Fault Tree / Event Tree approach (FTA)



Motivation and issues with current PSA softwares

- **Risk Informed Applications**

- ⌚ Evaluation and support of Plant / TechSpec Modifications
(e.g. relaxation)
- ⌚ Risk Informed In Service Inspection
- ⌚ Aging Programs
- ⌚ ILRT frequency relaxation
- ⌚ Outage optimization



Motivation and issues with current PSA softwares

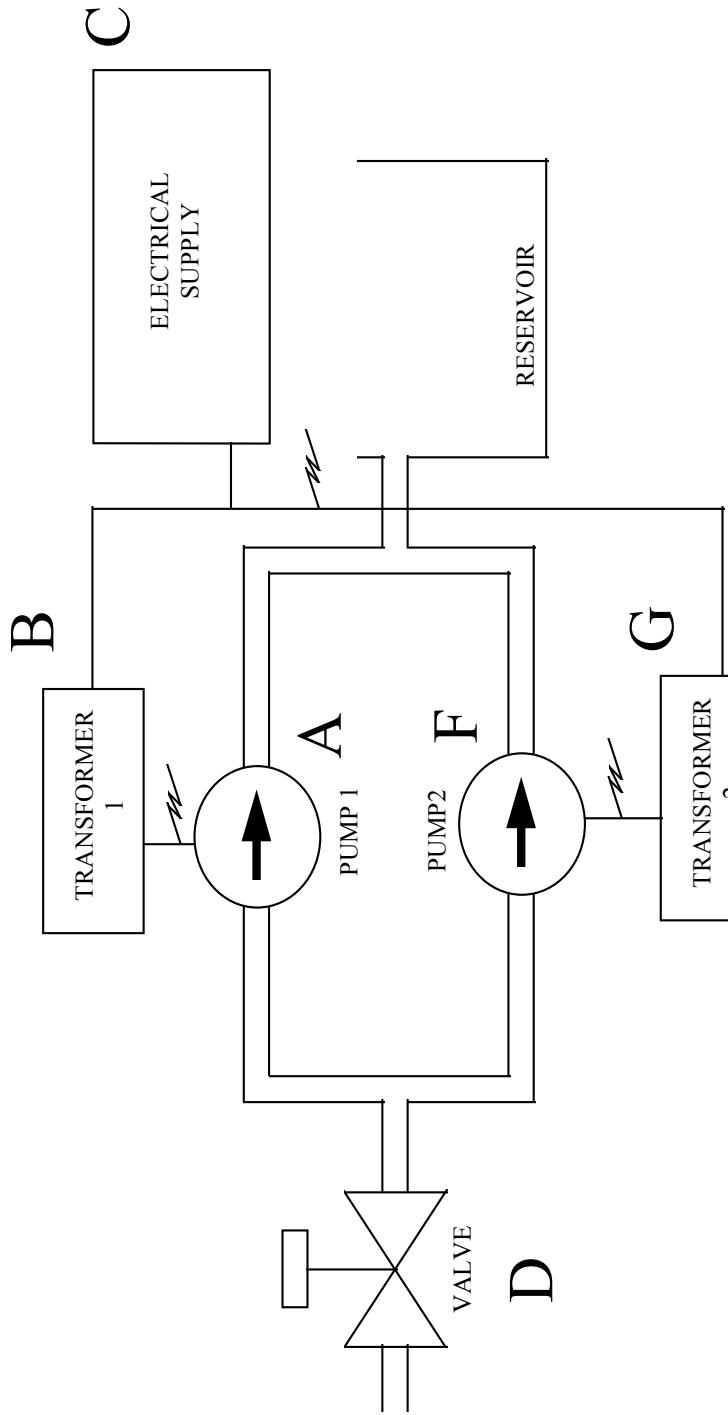
- Worldwide Fault Tree Analyses (FTA) are performed with codes that produce wrong results
 - ☼ Rare event approximation...
 - ☼ ... but not only !
- Resulting risk Importance measures are even more wrong (conservative or optimistic, no one can know)

$$RIF_{x_i} = \frac{CDF(x_i = 1)}{CDF} \nearrow$$



Motivation and issues with current PSA softwares

- What is the mean unreliability of the system's function, based on individual Basic Events probabilities ?

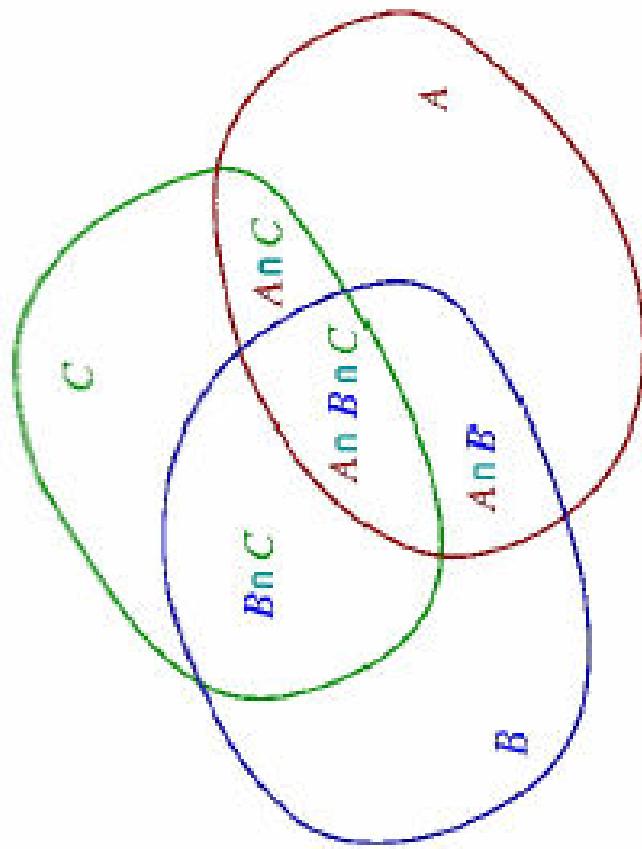


$$P_{top} = D + C + AF + AG + BF + BG$$

Motivation and issues with current PSA softwares

- The rare event approximation (Moivre's equation)

$$|A_1 \cup \dots \cup A_p| = \sum_{1 \leq i \leq p} |A_i| - \sum_{1 \leq i_1 < i_2 \leq p} |A_{i_1} \cap A_{i_2}| + \dots + (-1)^{p-1} |A_1 \cap \dots \cap A_p|$$

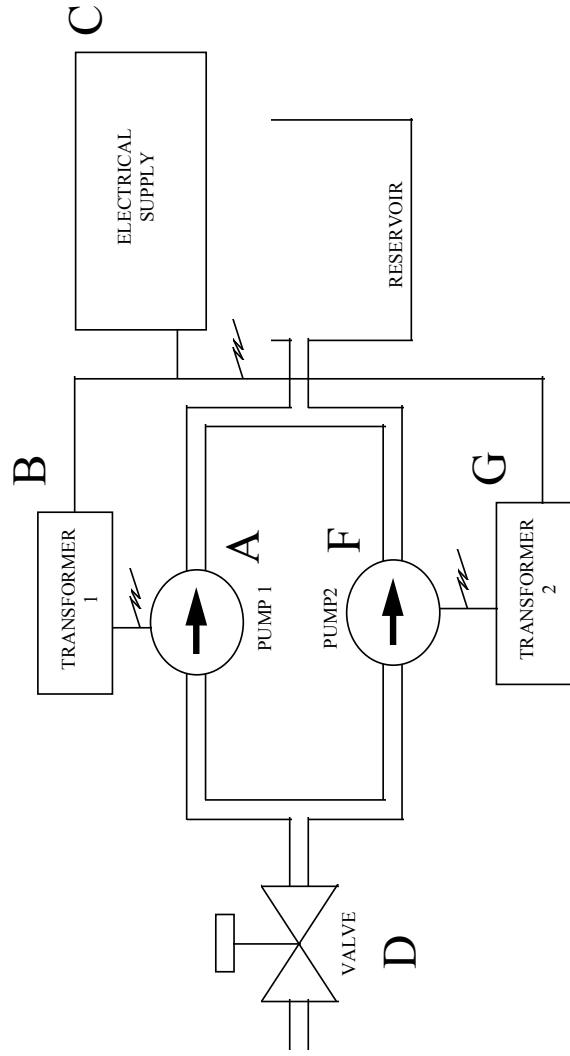


Motivation and issues with current PSA softwares

- Analytically correct result yields:

$$\begin{aligned} P_{\text{top}} &= (d + f + g + c - d f - d g - d c - f g - f c - g c + d f g + d f c + d g c + \\ &\quad f g c - d f g c) (a + b + c + d - a b - a c - a d - b c - b d - c d + \\ &\quad a b c + a b d + a c d + b c d - a b c d) \end{aligned}$$

$$\begin{aligned} &= [c - a (-1 + b) (-1 + c) (-1 + d) + b (-1 + c) (-1 + d) + d - c d] \cdot \\ &\quad [f - c (-1 + d) (-1 + f) (-1 + g) + d (-1 + f) (-1 + g) + g - f g] \end{aligned}$$

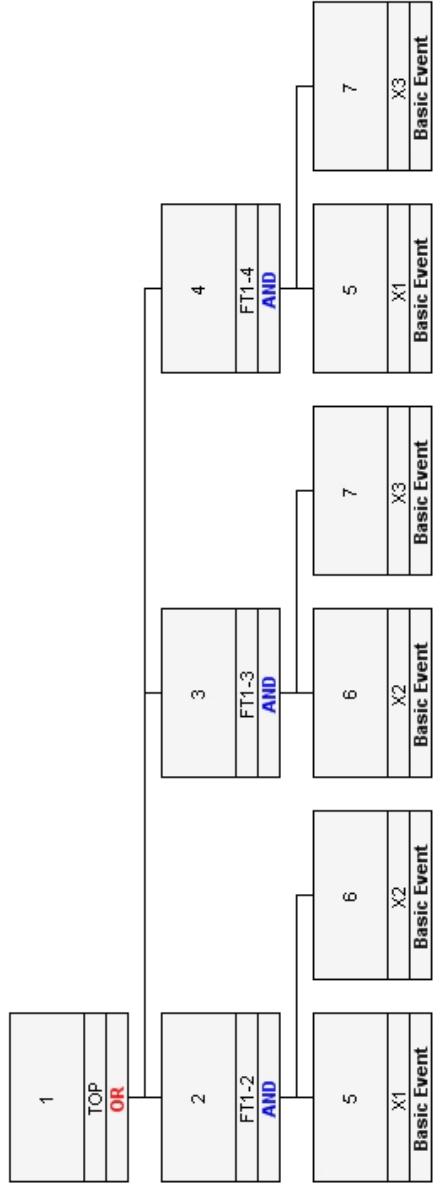


Motivation and issues with current PSA softwares

- Consider a « 2 out of 3 system » given by the Boolean equation:

$$T = (x_1 \wedge x_2) \vee (x_1 \wedge x_3) \vee (x_2 \wedge x_3)$$

$$p(x_1) = p(x_2) = p(x_3) := q$$



$$p(T) \neq p(x_1)p(x_2) + p(x_1)p(x_3) + p(x_2)p(x_3) = 3 \cdot q^2$$

$$p(T) = p(x_1)p(x_2) + p(x_1)\underbrace{(1-p(x_2))}_{\text{Success}}p(x_3) + \underbrace{(1-p(x_1))}_{\text{Success}}p(x_2)p(x_3) = 3q^2 - 2q^3$$



Motivation and issues with current PSA softwares

- Other issues include:
 - ✿ Wrong treatment of negative logic (e.g. forbidden maintenance unavailabilities according to TechSpec)
 - ✿ Quantification cutoff (typically 1E-12 to 1E-14)
 - ✿ Wrong interpretation of risk importance measures of components, systems and safety divisions (RIF, FV, etc.)
 - ✿ Treatment of exchange events
- Advanced PSA models include HRA, CCF, seismic and phenomenological events, where failure probabilities approach 1
- It is accepted that current quantification tools have reached their own limits [Rauzy, 2001]



Motivation and issues with current PSA softwares

- Develop a new PSA quantification methodology that:
 - ✿ Overcomes the deficiencies of the rare approximation, i.e. credit success branches, calculate the rare event up to infinite order
 - ✿ Yields a correct evaluation of Risk Importance Factors (RIFs)
 - ✿ Support the treatment of negative logic
 - ✿ Do not apply cutoff when generating the sequences
 - ✿ Improve calculation speed and result consistency



Binary Decision Diagrams (BDD) as an alternative

- Shannon expansion

$$x \rightarrow y_0, y_1 := (x \wedge y_0) \vee (\overline{x} \wedge y_1) := ite(x, y_0, y_1)$$

- Shannon expansion of t with respect to x

$$t = x \rightarrow t[1/x], t[0/x] \Rightarrow t = (x \wedge t[1/x]) \vee (\overline{x} \wedge t[0/x])$$

- ❖ $t[0/x]$ and $t[1/x]$ both contain one less variable than the expression t
- ❖ We can recursively find ITEs up to the basic elements 0 (false) and 1 (true)

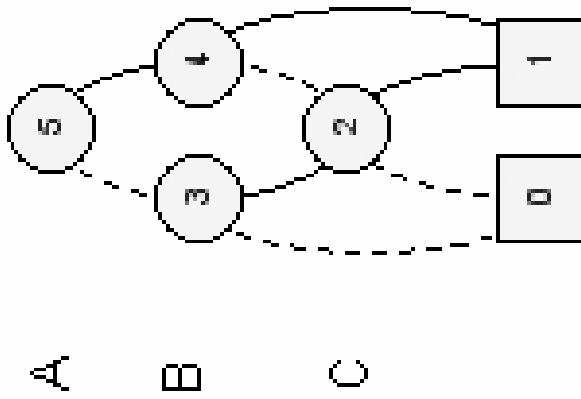


Binary Decision Diagrams (BDD) as an alternative

- Example for the „2 out of 3“ system $t = AB + BC + AC$

$t = A \rightarrow (t_0, t_1)$
⊗ $t_0 = B \rightarrow (0, t_{01})$
⊗ $t_1 = B \rightarrow (1, t_{10})$
⊗ $t_{01} = C \rightarrow (1, 0)$
⊗ $t_{10} = C \rightarrow (1, 0)$

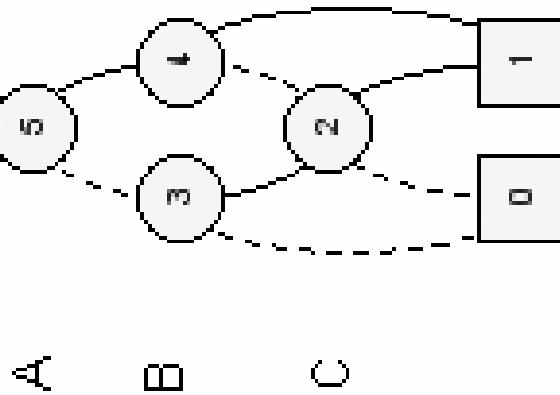
$\Rightarrow t = A \rightarrow (B \rightarrow (0, C \rightarrow (1, 0)), B \rightarrow (1, C \rightarrow (1, 0)))$



Binary Decision Diagrams (BDD) as an alternative

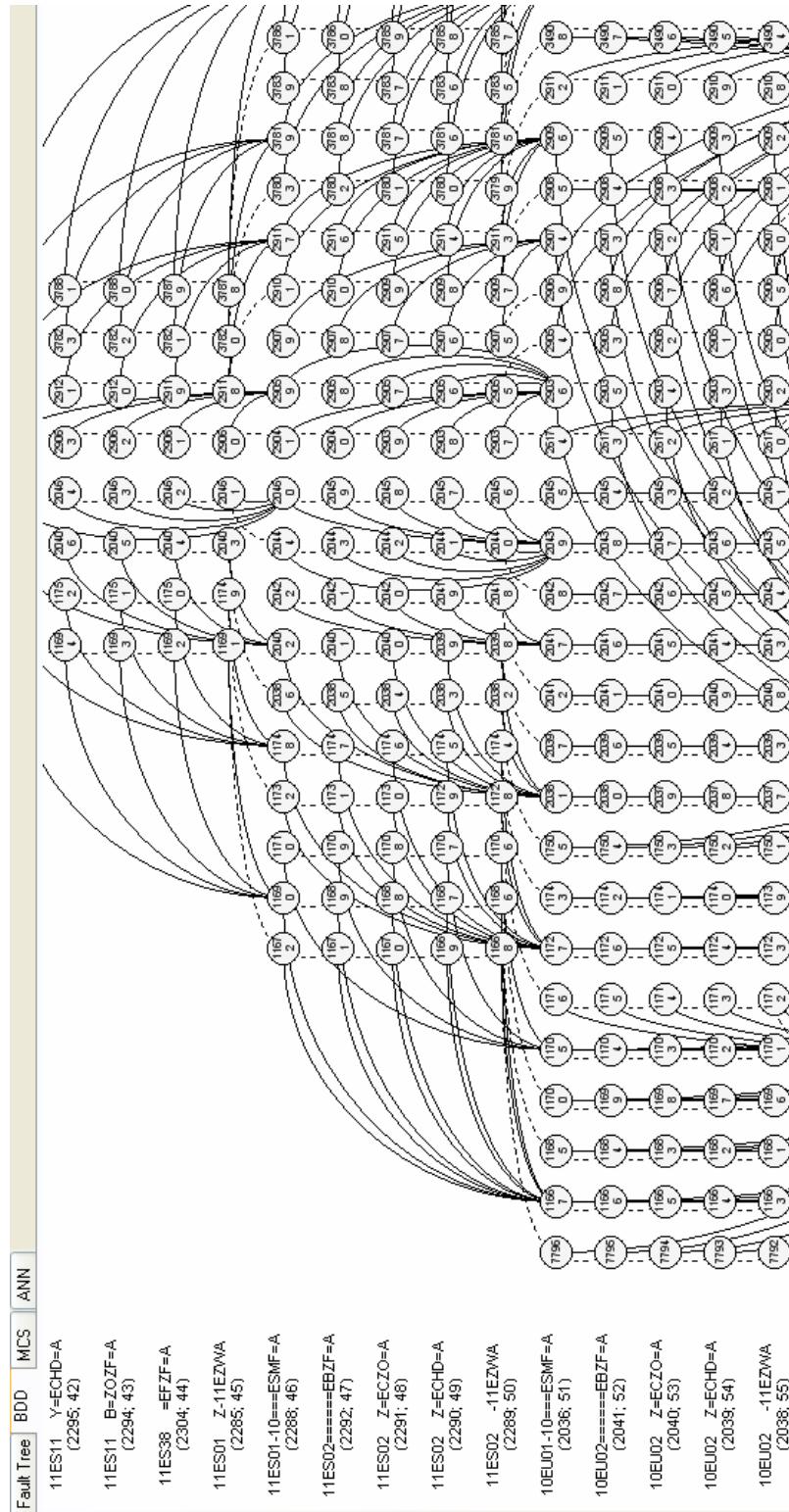
- Canonical formulation of Boolean equations!

- Example: $P_{t=true} = AB + A(1-B)C + (1-A)BC$



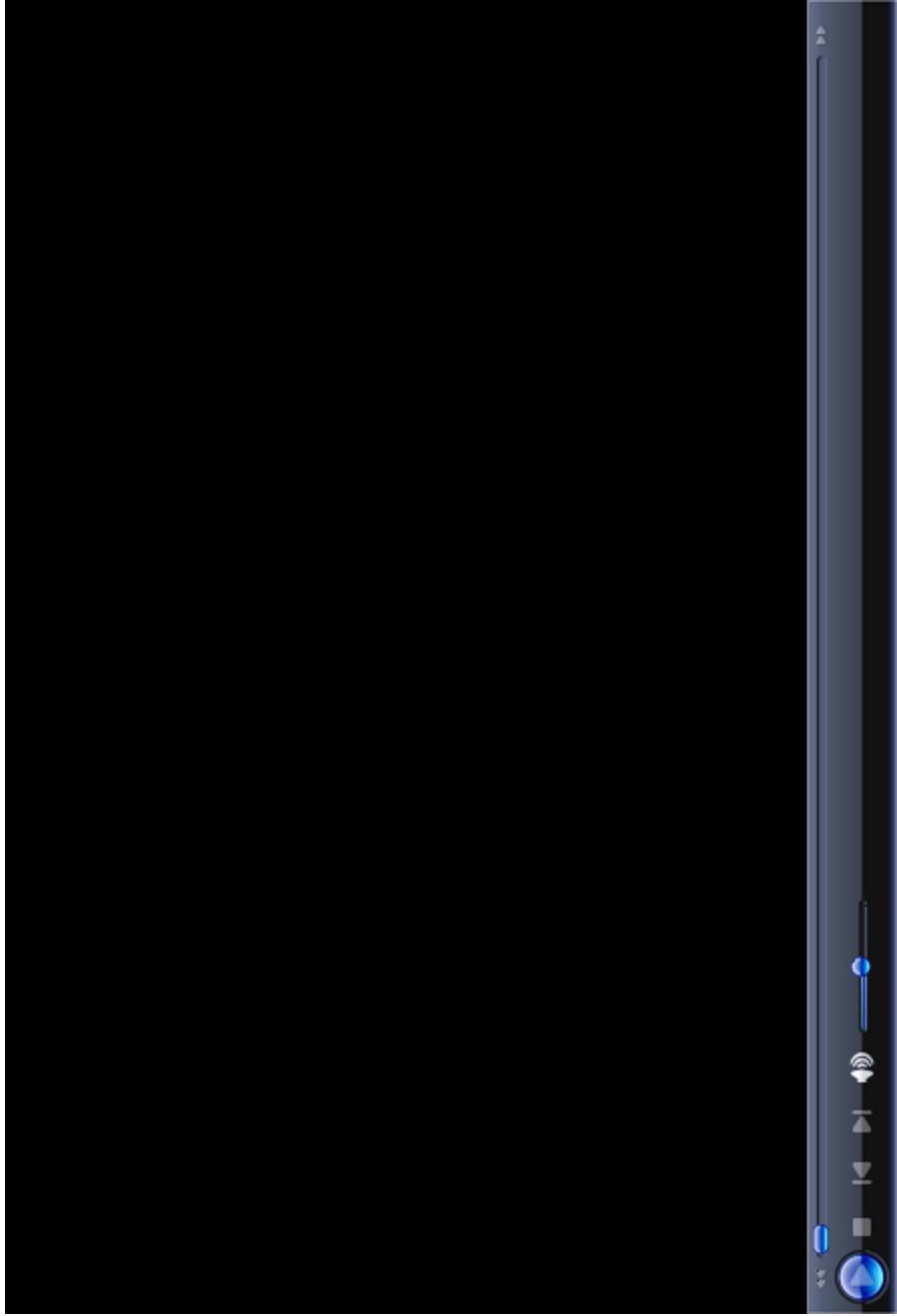
Binary Decision Diagrams (BDD) as an alternative

- BDD complexity is not related to the number of prime implicants of the encoded formula
- This small BDD (37620 nodes) encodes a total of 10^9 cutsets



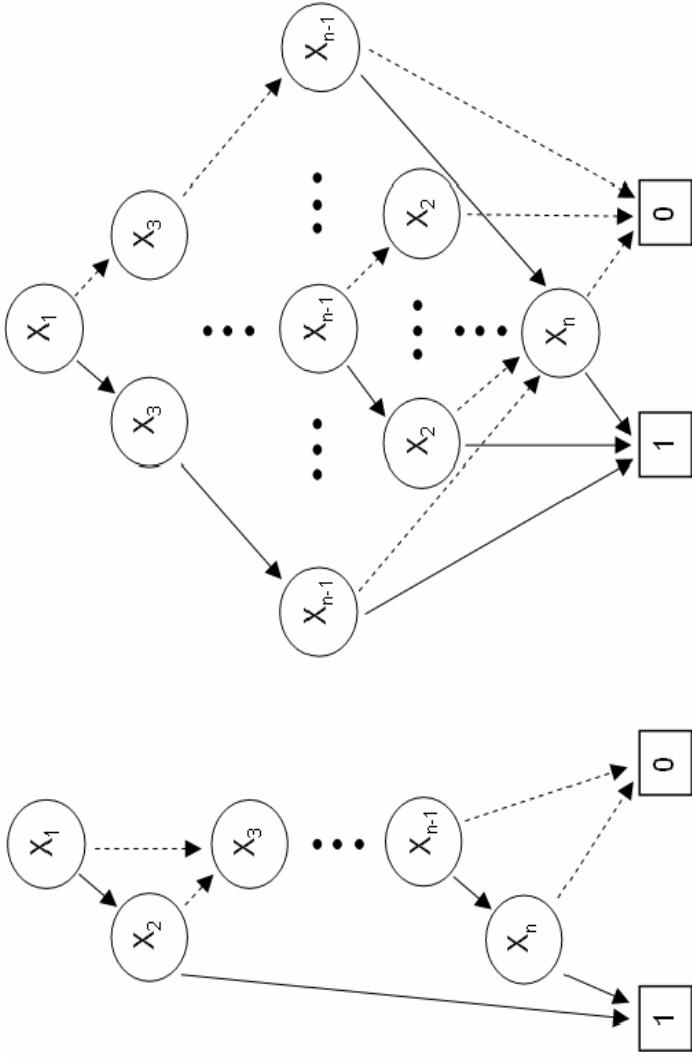
Binary Decision Diagrams (BDD) as an alternative

- Let's have a closer look at it !
- HPCS System of the Leibstadt NPP



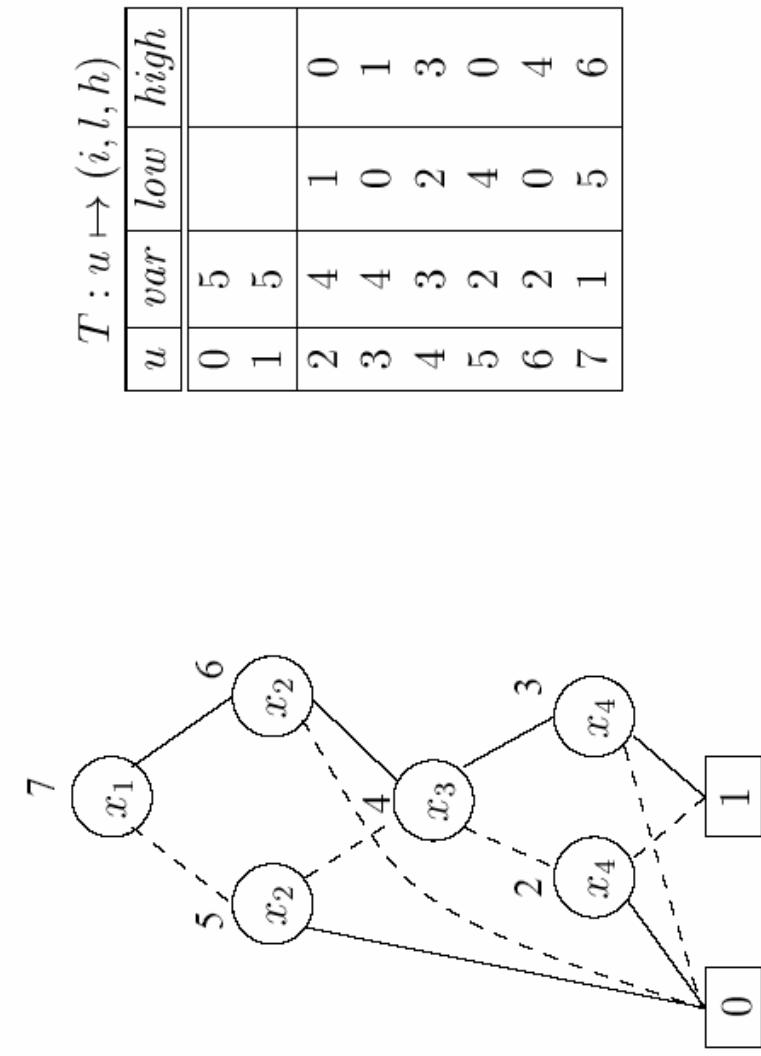
Binary Decision Diagrams (BDD) as an alternative

- Impact of variable order on BDD size
 - ✿ From linear ☺ to exponential ☹
 - ✿ ☹ Finding the best order is of **NP-Complete complexity** (Bollig / Wegener, 1996)

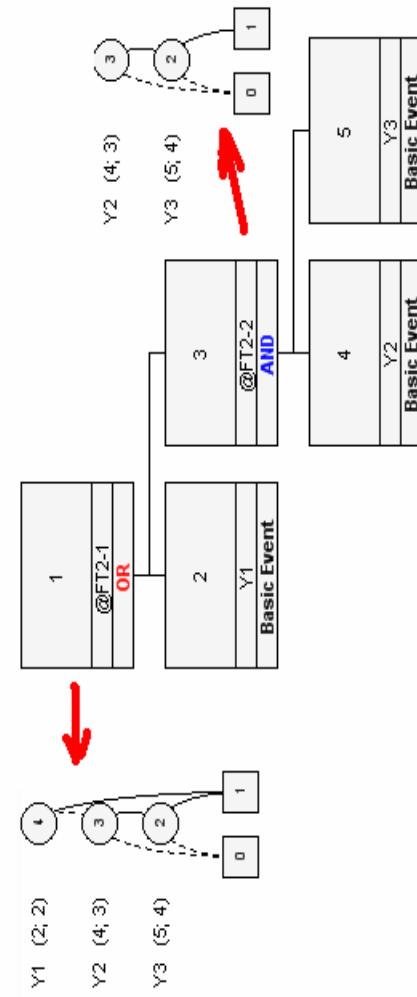
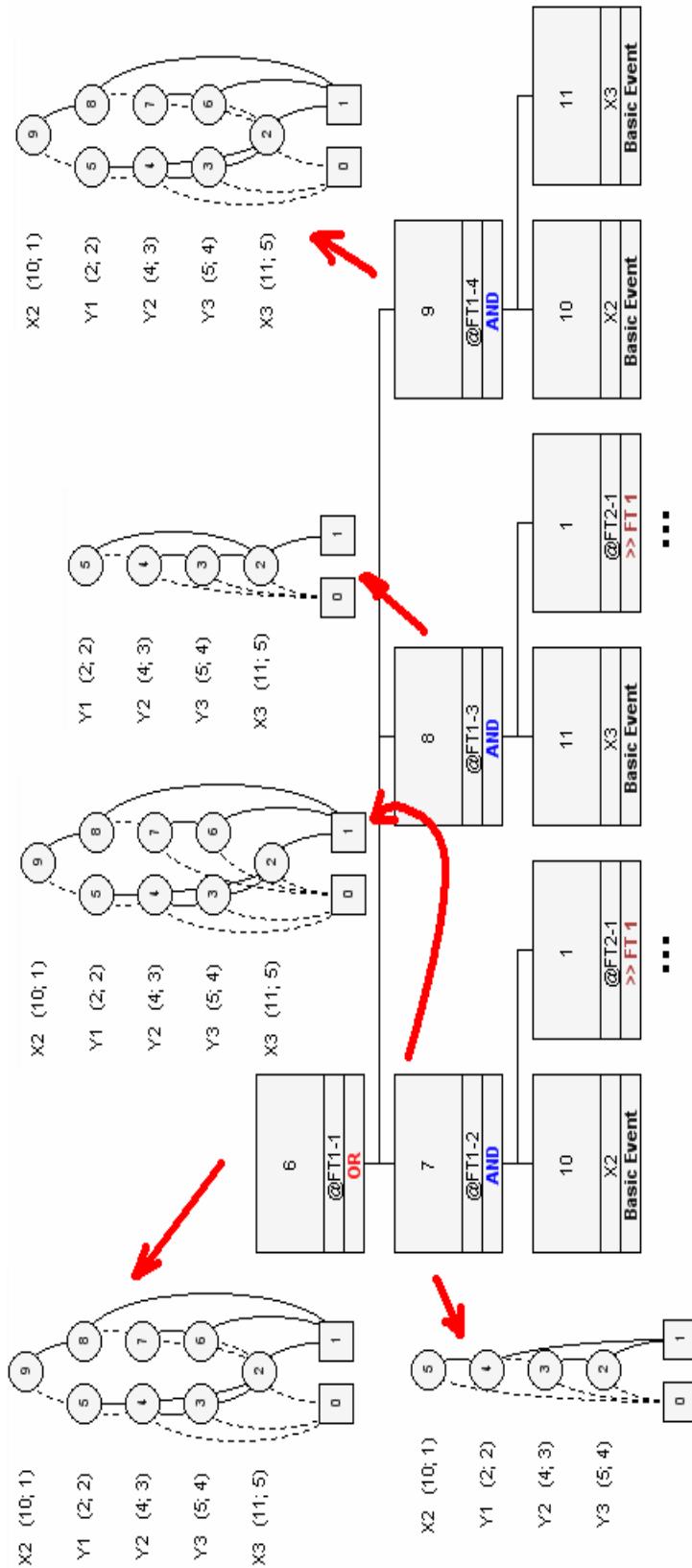


Binary Decision Diagrams (BDDs) as an alternative

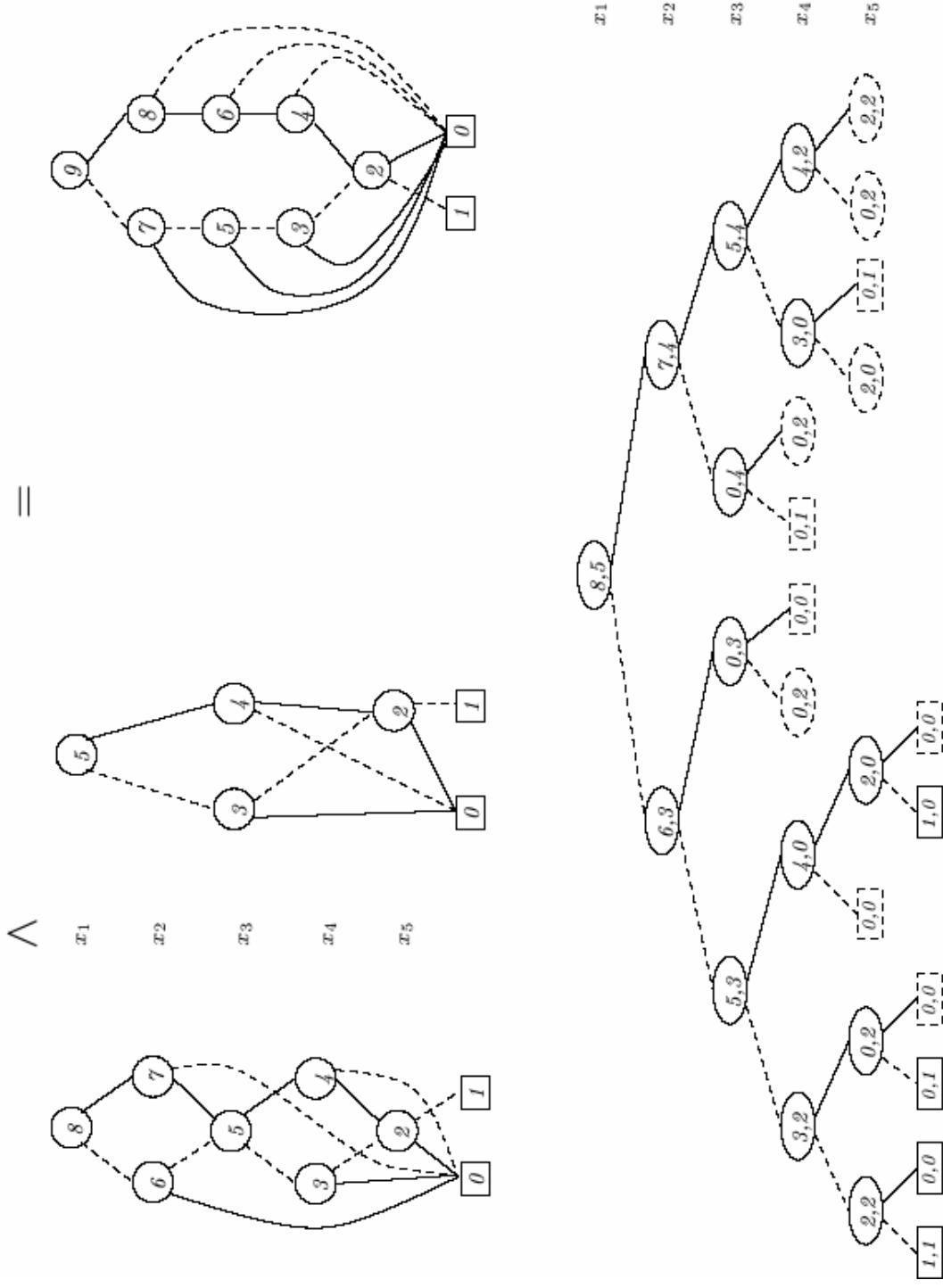
- Computer encoding of BDDs
 - ✿ Open-ended node table (dynamic)
 - ✿ Open-ended Hashtable (e.g. unique table)



Binary Decision Diagrams (BDD) as an alternative



Binary Decision Diagrams (BDD) as an alternative



Binary Decision Diagrams (BDD) as an alternative

Algorithm *Apply*[*T, H*](*op, u₁, u₂*)

Require: *u₁* and *u₂* the top nodes of the BDD to assemble.

Ensure: The resulting BDD.

```
1: if G(u1, u2)  $\neq \emptyset$  then
2:   return G(u1, u2)
3: else if u1  $\in \{0, 1\}$  and u2  $\in \{0, 1\}$  then
4:   u = op(u1, u2)
5:   else if var(u1) = var(u2) then
6:     u = newnode(var(u1), apply(low(u1), low(u2)), apply(high(u1), high(u2)))
7:   else if var(u1) < var(u2) then
8:     u = newnode(var(u1), apply(low(u1), u2), apply(high(u1), u2))
9:   else
10:    u = newnode(var(u2), apply(u1, low(u2)), apply(u1, high(u2)))
11: end if
12: G(u1, u2)  $\leftarrow$  u {Add to computation table}
13: return u {Returns node index}
```



Binary Decision Diagrams (BDD) as an alternative

- **Previous Work**

- BDD have been implemented in the early 90's for Integrated Circuits (IC) checking and IC optimization (16 and 32 bits)
- Some attempts to convert small to medium size models (typically with a few hundreds Basic Events) have succeeded
- All attempts with more Basic Events (>1000) have failed due to the exponential growth in complexity (BDD blow up)



Research and development at KKL and ETH Zurich

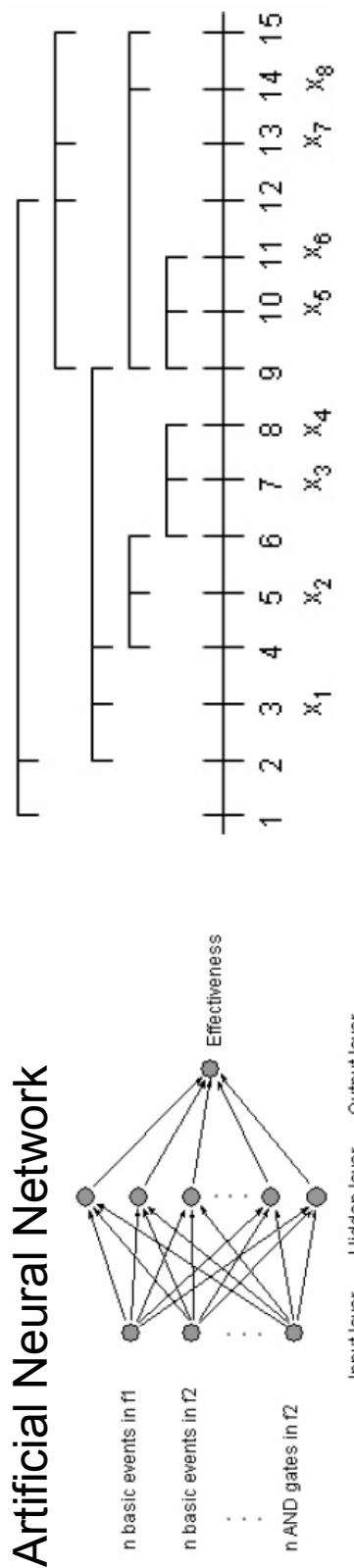
- Development of Fault Tree to BDD conversion engine
- Development of statistical measures
- Analysis of Fault Tree model pre-processing (rewriting)
techniques

❖ Basic Event occurrence based ordering

❖ Weights fan-out pre-processing → $W(v) = \begin{cases} 1 & \text{for basic events} \\ \sum_i W(v_i) & \text{for gates} \end{cases}$

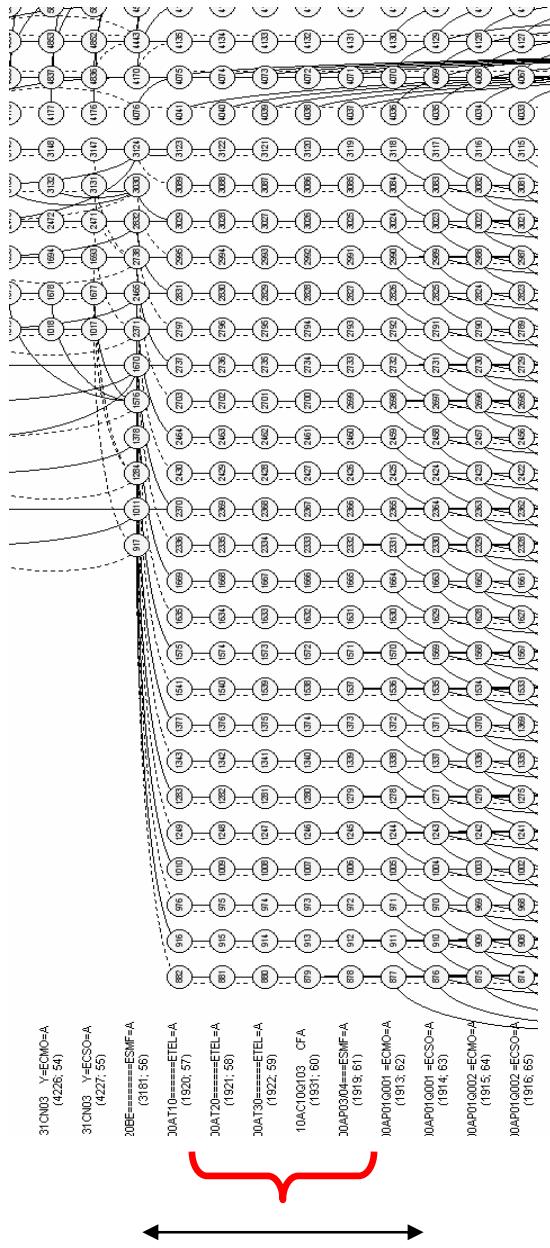
❖ Hypergraph optimization techniques

❖ Artificial Neural Network



Research and development at KKL and ETH Zurich

- Development of dynamic optimization techniques (e.g. Sifting, p-cut variable arrangement)
- Development of Group-Sifting for FTA



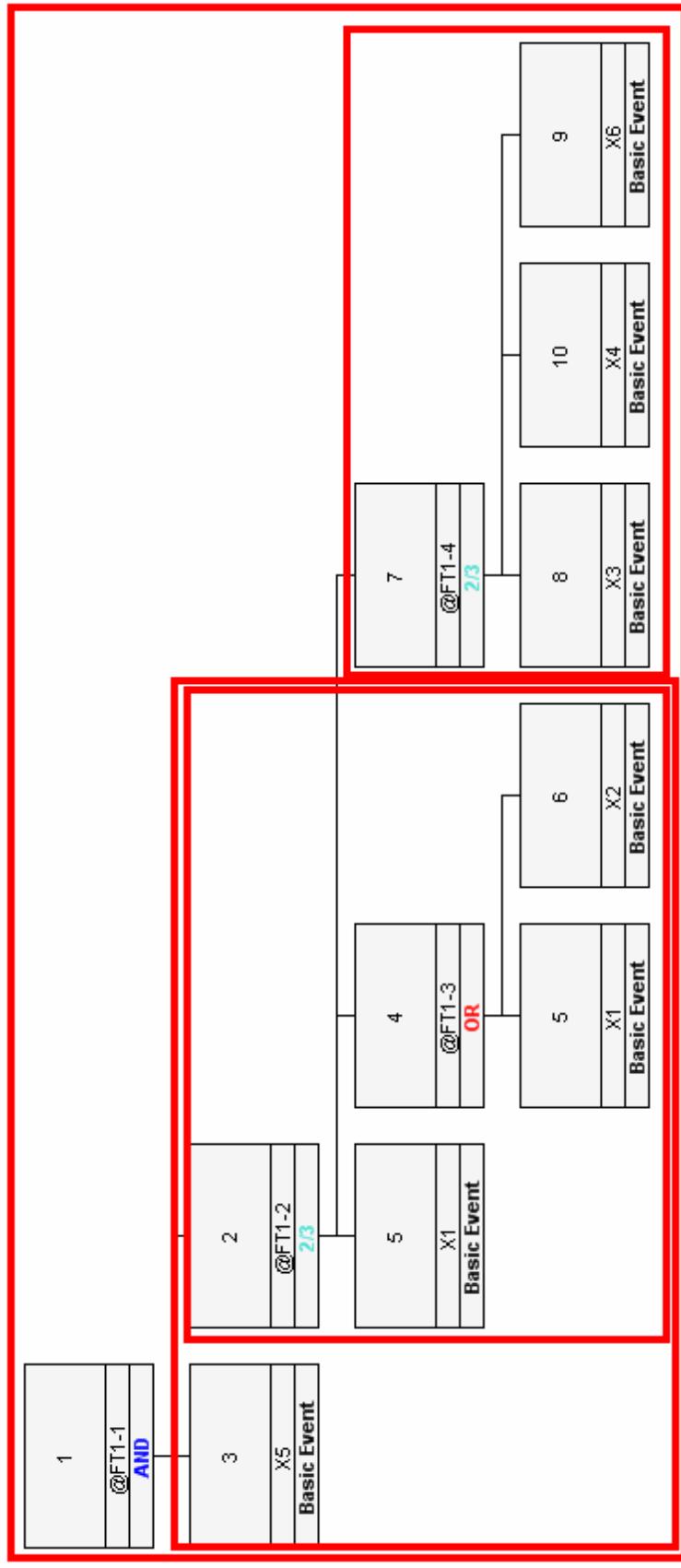
	DFLM	Regular sifting	Group-sifting
HPCSS	6545	3204	497
LPCSS	206'503	40'656	7763
RHR/A	306'339	99'945	11'948



Research and development at KKL and ETH Zurich

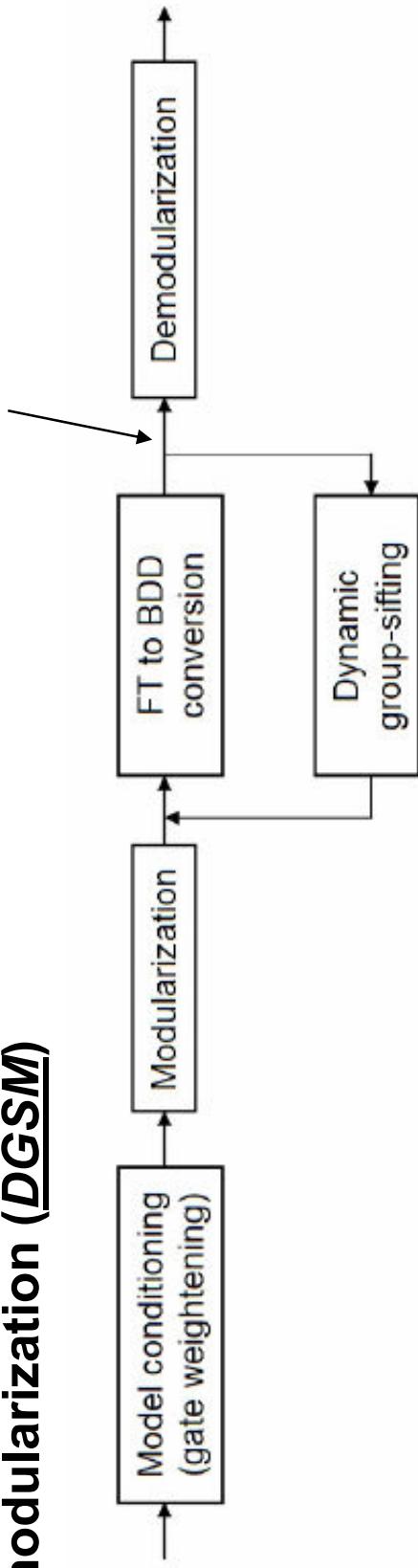
- Development and analysis of modularization techniques

✿ Occurrence vectors and detection criteria



Research and development at KKL and ETH Zurich

- Dynamic Group-Sifting using modularization (DGSM)

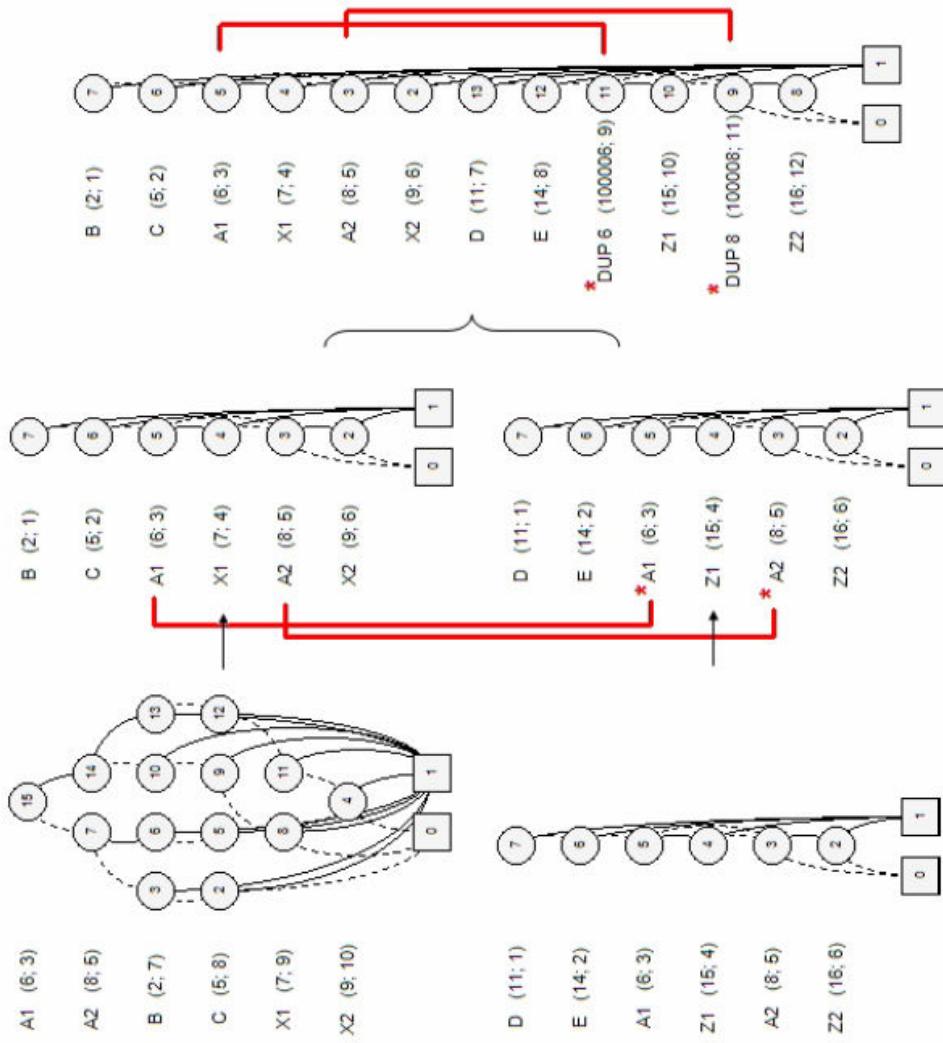


$$size_{BDD_R} > \left(\frac{size_{BDD_R}}{N} \right)^\alpha + 1 \times (size_{BDD_1} + size_{BDD_2})$$

Systems	DFLM	Dynamic group-sifting	Dynamic group-sifting using modularization
HPCS	6545	761	497
LPCS	206'503	3050	3053
RHR/A	306'339	3117	3117
RHR/B	impossible	52'447	21'177
ECCS	impossible	impossible	1'781'100 (CPU= 11 hours)

Research and development at KKL and ETH Zurich

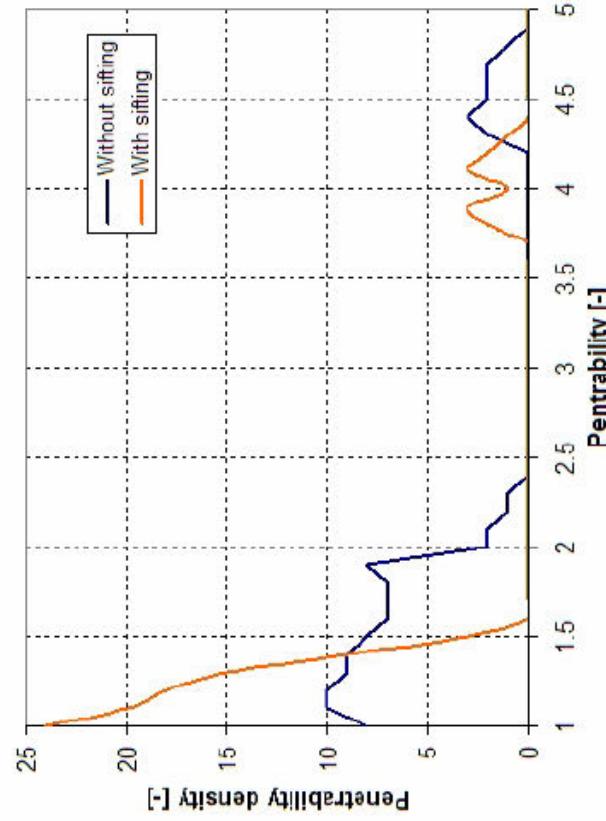
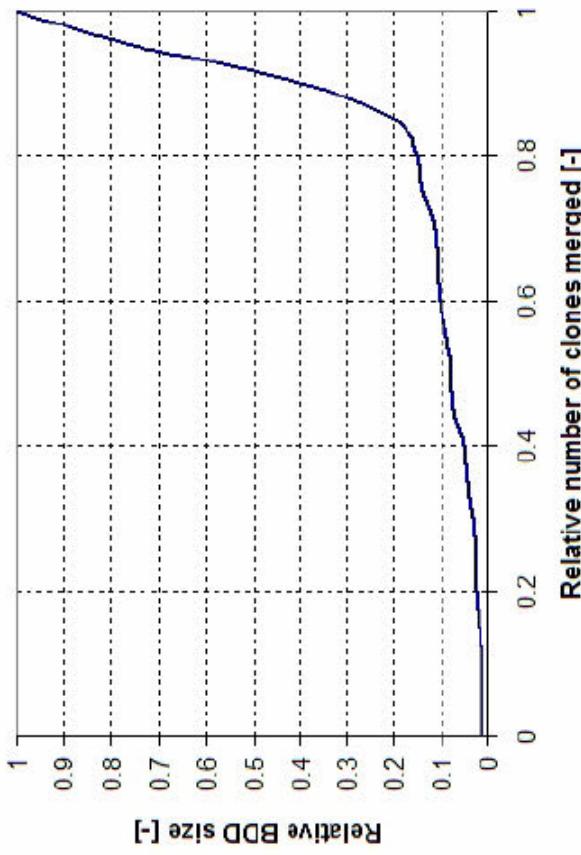
- Algorithm FUSION (BDD as objects)



Research and development at KKL and ETH Zurich

- **Performance of FUSION**

- ❖ About 90% of the variables can be merged without major impact on BDD size
- ❖ Penetrability ρ : effective identification of mathematical “hot spots”



$$p(x) = \frac{\text{Size}(\text{Opt}(BDD_{x \text{ merged}}))}{\text{Size}(BDD)}$$

Research and development at KKL and ETH Zurich

- **Improved Dynamic Group-Sifting Using Modularization (*IDGSM*)**
 - ☼ Further limitation of global perturbation when optimizing locally
 - ☼ Online identification and treatment of “hot spots” using penetrability spectrum
 - ☼ Improvement in the Group-Sifting algorithm
 - ☼ Use of genetic optimization algorithms
 - ☼ Generation and treatment of “clones” (Algorithm *FUSION*)



Insights and outlook

- **Results and insights:**

- FTA quantification using BDD requires complex algorithms and programming techniques
- The combination of global, static, dynamic and BDD objects techniques proved to be effective when dealing with large models
- We succeeded in converting the Leibstadt PSA model to a BDD form of more than 1'500'000 with 30 clones, for a total of about 3500 basic events
- The BDD covers a complete event tree sequence that includes reactor shutdown and reactor cooling with the eight Emergency Core Cooling Systems of the Leibstadt Nuclear Power Plant (including support systems)



Insights and outlook

- **The nuclear industry is facing a major issue (and is not yet fully aware of it):**
 - ✿ Worldwide probabilistic analyses are performed with codes that produce wrong results (conservative or optimistic, no one can know)
 - ✿ New IAEA requirements are impossible to address with existing FTA quantifiers (e.g. seismic assessments)
 - ✿ Utilities and authorities *still* trust the results of existing FTA quantifiers, applying approximations where they should not be applied
(→ blind-trust in numbers)
 - ✿ The techniques developed in this study raises the BDD approach to a mature technology for PSA model solving

